

## 2.8. Construction Trees Revisited (and Reversed)

Construction rules serve as a building code for formal sentences: just as city building codes set out the conditions for constructing a house that passes legal muster, so our construction rules state how to build a string of symbols that counts as a genuine, ‘legal’ sentence in the formal language (rather than just a string of formal gibberish).

But building codes can also be used to inspect a house that’s already constructed – say, as part of selling it. And our four construction rules can likewise be applied to assess a finished sentence whose construction we may not have witnessed. In that case we begin with the finished sentence, and hang its construction tree under it to prove it was constructed legally.

Recovering the construction tree in this way involves ‘un-building’ the sentence – **performing the construction process in reverse**. Since construction began with atoms and subsequently used three molecule-building rules, un-building uses those same **molecular rules in reverse** as **molecule-dissolving** procedures, leading back to the original atoms.

### Atomic Sentences:

1. Sentence letters are formal sentences.

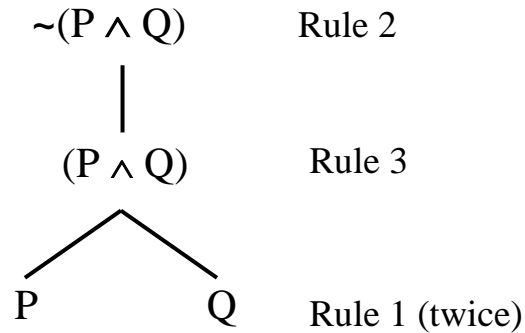
### Molecular Sentences:

2. If  $\blacktriangle$  is a formal sentence, then  $\sim\blacktriangle$  is a formal sentence.
3. If  $\bullet$  and  $\blacktriangle$  are formal sentences, then  $(\bullet \wedge \blacktriangle)$  is a formal sentence.
4. If  $\bullet$  and  $\blacktriangle$  are formal sentences, then  $(\bullet \vee \blacktriangle)$  is a formal sentence.

Specifically: since each molecular rule adds a connective (and, in the case of wedges and vels, a pair of parentheses), in reverse each molecular rule **removes a connective** (and with wedges and vels, a pair of parentheses).

The trick here is to decide where the un-building should begin – that is, which connective should be removed first. Let us call the last connective

added in the construction process the **main sign** of that sentence.<sup>1</sup> So in the following sentence the tilde is the main sign, as it was the last connective added in the construction of “ $\sim(P \wedge Q)$ ”.



The **left-most symbol** turns out to be a reliable clue as to which rule applied last in construction (and hence which connective is the main sign of the sentence). For the output of Construction Rule 2 has a tilde as left-most symbol; whereas Construction Rules 3 and 4 leave a left parenthesis as left-most symbol.

2. If  $\blacktriangle$  is a formal sentence, then  $\sim\blacktriangle$  is a formal sentence.
3. If  $\bullet$  and  $\blacktriangle$  are formal sentences, then  $(\bullet \wedge \blacktriangle)$  is a formal sentence.
4. If  $\bullet$  and  $\blacktriangle$  are formal sentences, then  $(\bullet \vee \blacktriangle)$  is a formal sentence.

The above sentence “ $\sim(P \wedge Q)$ ” has a tilde as left-most symbol; and that alone tells us it’s a negation, the output of Rule 2. Reading its construction tree from top to bottom outlines its un-building: Rule 2 in reverse removes a tilde, yielding “ $(P \wedge Q)$ ,” a conjunction produced by Rule 3; and Rule 3 in reverse removes a wedge and parentheses, leaving the two sentence letters “P” and “Q”. (Sentence letters can’t be un-built with any molecular rule in reverse, since they have no connectives to remove.)

<sup>1</sup> Following (Partee, ter Meulen, and Wall 1990: XX). Calling it the “main sign” rather than “main connective” allows the term to be extended to quantifiers in Chapters Five and Six.

Consider next a sentence without a construction tree.

$$((P \wedge Q) \vee \sim R)$$

Being a molecular sentence, “ $((P \wedge Q) \vee \sim R)$ ” must have been the output of one of the three molecular rules. And since the left-most symbol here is a parenthesis, it could only be the output of Rule 3 or Rule 4.

As a matter of fact this sentence is a **disjunction**, the product of Rule 4.<sup>2</sup>

4. If  $\bullet$  and  $\blacktriangle$  are formal sentences, then  $(\bullet \vee \blacktriangle)$  is a formal sentence.

Its main sign is thus a vel – the very connective that brought those parentheses with it.

$$\underline{( (P \wedge Q) \vee \sim R )}$$

From this output we work back to the two inputs, by applying Rule 4 in reverse. Since Rule 4 adds a vel and outer parentheses, Rule 4 in reverse **removes** a vel and outer parentheses.

**Rule 4 in Reverse:** remove a vel, and the outermost pair of parentheses.

This leads us back to the two sentences being linked together by the vel.

$$\begin{array}{c} ((P \wedge Q) \vee \sim R) \\ \diagdown \quad \diagup \\ (P \wedge Q) \quad \sim R \end{array}$$

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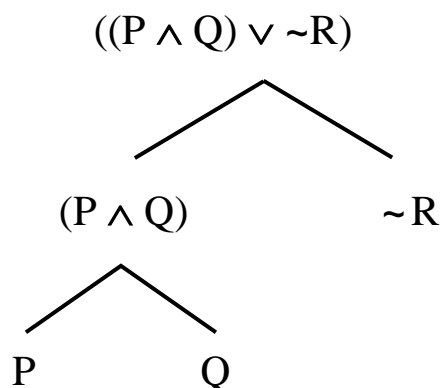
<sup>2</sup> A procedure for showing mechanically *why* the vel, and not the wedge, is the main connective here is addressed in 2.8.1. Problem C.

The left part, “ $(P \wedge Q)$ ,” is a smaller molecule with a wedge as its main sign. “ $(P \wedge Q)$ ” is the product of Rule 3, the conjunction rule.

3. If  $\bullet$  and  $\blacktriangle$  are formal sentences, then  $(\bullet \wedge \blacktriangle)$  is a formal sentence.

Rule 3 in reverse **removes** a wedge and outer parentheses.

**Rule 3, in Reverse:** remove the outermost pair of parentheses, and take a conjunction sign from between the two parts.



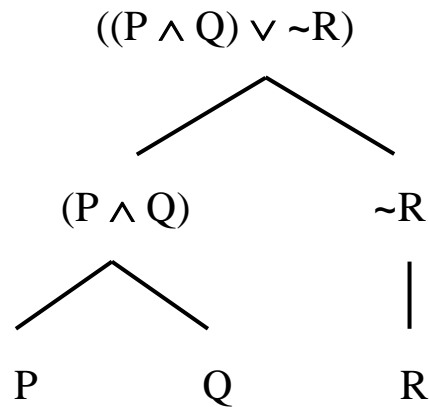
Since “P” and “Q” are atoms, they cannot be un-built by any molecular rule in reverse.

But “~R,” on the right of the tree, is a molecule susceptible to disassembly. “~R” has a tilde as its left-most symbol – meaning it’s a negation, built by Rule 2.

2. If  $\blacktriangle$  is a formal sentence, then  $\sim\blacktriangle$  is a formal sentence.

Rule 2 in Reverse removes a tilde from the left of the sentence.

**Rule 2 In Reverse:** remove a tilde from the left of the sentence.



This illustrates the general strategy for recovering the construction tree for any formal sentence: break down the sentence using the three molecular rules in reverse, until only atomic sentences (sentence letters) remain.

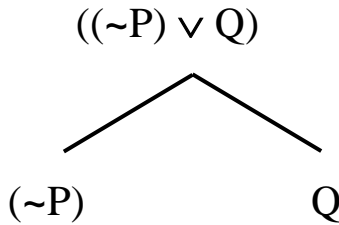
It turns out that **any genuine (legal) formal sentence can be un-built, by the molecular rules in reverse, to just sentence letters.**<sup>3</sup> And anything which isn't a formal sentence can't be un-built back to sentence letters in this way.<sup>4</sup>

<sup>3</sup> In fact it can be shown that a legal formal sentence will have **one and only one** (rule-following) construction tree.

<sup>4</sup> In practice – provided the sentence isn't overwhelmingly large – we can usually just 'look and see' which connective is the main sign, and so how sentence unbuilding should go for each step. But for a computer that applied the molecular rules blindly in search of a proper construction tree (without first determining which connective is the main sign) the rule would instead be: a legal formal sentence will yield a construction tree (adhering to the construction rules) with only sentence letters at the bottom – along with however many bogus 'trees' without sentence letters at the bottom. (For instance, by using Rule 3 in reverse to remove a wedge and outermost parentheses, we could un-build the perfectly fine formal sentence " $((P \wedge Q) \vee R)$ " into " $(P)$ " and " $(Q \vee R)$ " – neither a sentence letter, and neither susceptible of further un-building.) On that blind 'exhaustive search' approach, anything which isn't a legal formal sentence will yield **not even one** (rule-following) tree that resolves into sentence letters.

The steps we might use to teach the computer to apply the reverse construction procedure less blindly – yielding just the one correct tree for a formal sentence, without wading through a jungle of pseudo-trees – is addressed in 2.18.1, Problem C.

For instance, the following string of symbols is a piece of formal gibberish, as a reverse construction tree shows.



By removing a wedge and outer parentheses, Rule 3 in reverse can certainly un-build “ $((\sim P) \vee Q)$ ” into parts “ $(\sim P)$ ” and “ $Q$ ”. But **no molecular rule in reverse can un-build “ $(\sim P)$ ”**. Because the left-most symbol is a left parenthesis, Rule 2 can’t apply. Rule 3 in reverse removes a pair of parentheses and a wedge – meaning it can’t apply to “ $(\sim P)$ ,” which has no wedge to remove. For the same reason Rule 4 can’t be used here, because in reverse Rule 4 removes parentheses and a vel. Since molecular rules in reverse can’t un-build it entirely to sentence letters, “ $((\sim P) \vee Q)$ ” is shown not to be a legal formal sentence, but mere symbolic gibberish.

Reverse construction trees have a variety of applications. Separating the legal formal sentences from their illegal imposters is one, as we’ve seen. However, recovering the construction tree for a formal sentence will later also prove essential to the truth table test of validity.

But we stress finally that understanding how a formal sentence is constructed also makes clear that sentence’s **logical meaning** – whether, for instance, “ $\sim(P \wedge Q)$ ” is denying a conjunction, or asserting one (with a denial as one of its parts).<sup>5</sup> For that is crucial to proper translation from English to the formal language – a topic to which we now return.

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<sup>5</sup> It’s the **denial** of a conjunction – because its main sign is a **tilde**.